

## 11

## Derived Schemes.

Before discussing deformation theory we will introduce a class of prestacks that are natural generalizations of classical schemes. The philosophy here is that derived schemes are prestacks w/ certain properties, yet we don't need to introduce any extra structures, for instance what a locally ringed  $\infty$ -topos is.

Def'n: A prestack  $\mathcal{Z}$  is said to be a scheme if

- (i)  $\mathcal{Z}$  satisfies étale descent;
- (ii)  $\mathcal{Z} \rightarrow \mathcal{Z} \times \mathcal{Z}$  is affine schematic, and  $\forall \overline{T} \rightarrow \mathcal{Z} \times \mathcal{Z}$  aff.

the map  $\mathcal{C}_{\overline{T} \times \mathcal{Z}}^{\mathcal{Z} \times \mathcal{Z}} \rightarrow \mathcal{C}_{\overline{T}}$  is a closed embedding;

- (iii)  $\exists$  a collection  $\{f_i: S_i \rightarrow \mathcal{Z}\}_I$  (Zariski atlas) s.t.
  - $f_i$  is an open embedding,
  - $\coprod_I \mathcal{C}_{\overline{T} \times S_i}^{\mathcal{Z}} \rightarrow \mathcal{C}_{\overline{T}}^{\mathcal{Z}}$  cover  $\mathcal{C}_{\overline{T}}^{\mathcal{Z}} \quad \forall \overline{T} \rightarrow \mathcal{Z}$  aff.

We let  $\text{Sch} \hookrightarrow \text{Stk} \hookrightarrow \text{PStk}$  denote the subcategory of schemes.

Rk. 1. (i)'  $\mathcal{Z}$  satisfy Zariski descent  $\Rightarrow$  (i)  $\mathcal{Z}$  satisfy étale descent  $\Rightarrow$  (i)''  $\mathcal{Z}$  satisfy flat descent.

2. We are restricting to the case of ~~separated~~ separated in this definition for ease of exposition. The non separated cases will be included in  $n$ -geometric stacks to be defined later.

We discuss some formal properties of the definition.

For  $Z$  a scheme and  $\{f_i: S_i \rightarrow Z\}$  a Zariski atlas, one has that

$\coprod_I S_i \rightarrow Z$  is an étale surjection.

$$\Rightarrow \bigotimes_{\mathbb{Z}} L\left(\coprod_I S_i / Z\right)_{\text{pstk}} \xrightarrow{\cong} Z. \quad (\star)$$

Lemma:

(Lifting ~~properties~~ open embeddings / étale morphisms.) Let  $f: Z' \rightarrow Z$  be an affine Zariski or étale morphism between schemes. For  $T \rightarrow Z$ , where  $T \in \text{Sch}^{\text{aff}}$ , one has an equivalence (of spaces).

$$\left\{ \begin{array}{c} Z' \\ \swarrow \quad \downarrow \\ T \rightarrow Z \end{array} \right\} \cong \left\{ \begin{array}{c} \mathbb{A}^1_{Z'} \\ \swarrow \quad \downarrow \\ \mathbb{A}^1_T \rightarrow \mathbb{A}^1_Z \end{array} \right\}.$$

Pf: Indeed, for  $Z' \rightarrow Z$  in  $\text{Sch}^{\text{aff}}$  this follows from the equivalence of the (small) étale site over  $S$  &  $\mathbb{A}^1_S$ . For schemes this follows from using the presheaf  $(\star)$ .

The ~~new~~ presentation of  $(\star)$  is useful to determine when one has a scheme for a given underlying classical scheme.

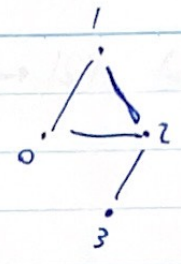
Digression: given an  $\infty$ -category  $\mathcal{L}$  a groupoid object in  $\mathcal{L}$  is a functor.

$$X: \Delta^{\text{op}} \rightarrow \mathcal{L} \quad \text{s.t.} \quad \forall n \geq 0 \quad \& \text{ a partition } [n] = S \cup S' \text{ w/ } S \cap S' = \emptyset \text{ one has.}$$

$$X([n]) \xrightarrow{\cong} \underset{X(\emptyset)}{X(S)} \times_{X(S')} X([n]) \quad \text{in } \mathcal{L}.$$

$$\text{E.g.: } - X([0]) \xrightarrow{\cong} \bigotimes_{\mathbb{Z}} X([1]) \times_S X([1]).$$

$$\begin{aligned}
 - X(\mathbb{A}^1) &\xrightarrow{\cong} X(\mathbb{A}^1) \times X(\mathbb{A}^1) \\
 &\quad \quad \quad X(\mathbb{A}^1) \\
 \Rightarrow X(\mathbb{A}^1) &\rightarrow X(\mathbb{A}^1) \times X(\mathbb{A}^1) \times X(\mathbb{A}^1) \\
 &\quad \quad \quad X(\mathbb{A}^1) \quad X(\mathbb{A}^1)
 \end{aligned}$$



Fact: In  $\mathcal{S} = \mathcal{S}_{pc}$ , or more generally  $\mathcal{S}$  an  $\infty$ -topos, e.g.  $\mathcal{S} = \text{PStk}$  or  $\text{Stk}$ , any groupoid object  $X_0: \mathbb{A}^{\text{op}} \rightarrow \mathcal{S}$  is effective, i.e.

$$X \cong (X_0 / X_{-1}) \quad \text{where } X_{-1} := \text{colim}_{i \in \mathbb{A}^{\text{op}}} X_i \text{ sub.}$$

$(X_0 / X_{-1})$  is the Čech nerve, i.e.  $(\dots \rightrightarrows X_0 \xrightarrow{X_0} X_0 \rightrightarrows X_0)$  in a topos.

This result should be compared w/ the fact in ordinary category th. that any equivalence relation is effective, i.e.  $R \subseteq X \times X$  an equivalence relation then

$$\begin{aligned}
 p: R \rightarrow X \times X \text{ is s.t. } R \cong X \times X \\
 \forall S \in \mathcal{C} \quad \text{Hom}(S, R) \cong \text{Hom}(S, X) \times \text{Hom}(S, X) \\
 \cong \text{is an eq. rel.}
 \end{aligned}$$

Let  $S^\bullet: \mathbb{A}^{\text{op}} \rightarrow \text{PStk}$  be a groupoid object in  $\text{PStk}$ .

Consider  $Z := L(|S^\bullet|)$ .

Prop: If  $S^\bullet = \coprod_I S_i^\bullet$ ,  $S_i^\bullet = \coprod_J S_j^i$  are disjoint unions of affine schemes &

each  $S_j^i \rightarrow S_i^\bullet$  is an open embedding.  
 $\mathcal{C} \in Z$  is a scheme &  $\coprod_I \mathcal{C} \times S_i^\bullet \rightarrow \mathcal{C} \times Z$  is a Zariski atlas.

Then: (i)  $Z$  is a scheme;  
 (ii)  $\coprod_I S_i^\bullet \rightarrow Z$  is a Zariski atlas.

Idea: (a) It is enough to show that each  $S_i^\bullet \rightarrow Z$  is an affine open embedding. (Exercise.)  
 (b)  $X \rightarrow Y$  in  $\text{PStk}$  is affine schematic  $\Leftrightarrow L(X) \rightarrow L(Y)$  is affine schematic (Consequence of the fact that affine schemes are sheaves.)  
 (c) Enough to check  $S_i^\bullet \rightarrow |S^\bullet|$  is an aff. open embedding.



We claim  $Z_0$  is a classical scheme. In particular,  $Z_0$  is 0-truncated.

Proof: ~~Let  $T_0 \in \text{Sch}^{\text{aff}}$  we claim  $Z_0(T_0)$  is  $k$ -truncated  $\Leftrightarrow T_0 \times_{Z_0} Z_0$  is  $k$ -truncated.~~

Notice that  $Z = (|S'|_{\text{pshtk}})$  is 0-truncated if  $|S'|_{\text{pshtk}}$  is 0-truncated for  $T_0 \in \text{Sch}^{\text{aff}}$

$$|S'|_{\text{pshtk}}(T_0) = |S'(T_0)|_{\text{Spec}}$$

$S'(T_0)$  satisfies the Kan condition, i.e.

$S'(T_0)(\Delta^n) \rightarrow S'(T_0)(\Delta^i)$  induces a surjection on  $\pi_0$ .

(It is actually an isom. since  $S'(T_0)$  is a groupoid object in  $\text{Spec}$ .)

Thus,

$$\pi_0 |S'(T_0)|_{\text{Spec}} \cong \pi_0 (\text{Fib}(S'(T_0)(\Delta^n) \rightarrow S'(T_0)(\Delta^i)))$$

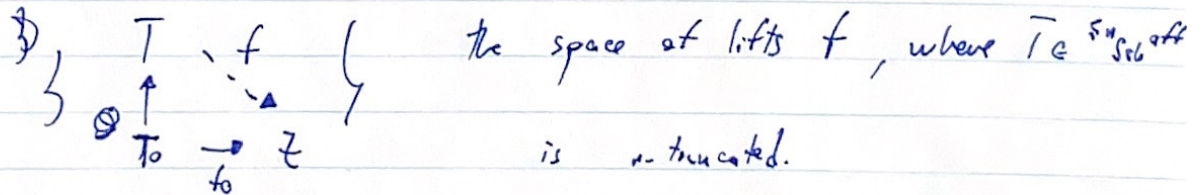
In particular, one has  $\pi_0 = \pi_0^{|\text{Spec}|} (\text{Fib}(\text{Hom}_{\text{Sch}^{\text{aff}}}(T_0, S') \rightarrow \text{Hom}_{\text{Sch}^{\text{aff}}}(T_0, S^0) \times \text{Hom}_{\text{Sch}^{\text{aff}}}(T_0, S^1)))$

but  $\alpha$  is injective since  $S^1 \cong S^0 \times S^0$  &  $Z \rightarrow Z \times Z$  affine closed imply that  $S^1 \rightarrow S^0 \times S^0$  is locally closed.

Similarly,  $\pi_0(|S'(T_0)|_{\text{Spec}}) = 0 \quad \forall k \geq 2$ .

for  $Z \in \text{Sch}^{\text{pshtk}} \cap \text{Sch} =: \text{Sch}^{\text{aff}}$

~~More generally~~ Further more, one can prove that given  $T_0 \in \text{Sch}^{\text{aff}}$



This gives  $\text{Hom}_{\text{Sch}^{\text{aff}}}(T_0, Z) \in \text{Sch}^{\text{aff}}$  is  $n$ -truncated. □

As with stacks one says a scheme  $Z \in \text{Sch}$  is  $n$ -cocommutative, if it is  $n$ -connective as a stack, i.e.

$$\text{Hom}_{\text{Sch}^{\text{aff}}}(\text{Hom}_{\text{Sch}^{\text{aff}}}(Z, Z), Z) \xrightarrow{\cong} Z$$

We summarize <sup>some</sup> results related to this:

Cor:  $Z \in \text{Sch}$  is  $n$ -coconnective iff ~~iff~~ <sup>one of</sup>  $\exists S^0 \rightarrow Z$  a Zarski cover  
w/  $S^0 \leftarrow \coprod_I S_i^0$ , w/  $S_i^0 \in \text{Sch}^{\text{aff}}$ .  
<sup>↳  $\text{Lk} \text{Sch}^{\text{aff}} \rightarrow \text{Sch}^{\text{aff}}$  identifies  $\text{Sch}^{\text{aff}}$  w/  $n$ -coconnective schemes.</sup>

We also mention the relation between the above condition & convergence.

(See Gr-I, Chapter 2 §3.4).

Prop: For  $Z \in \text{Sch}$ ,  $Z$  is ~~co~~ convergent.

Also given  $Z \in \text{Sch}$  a convergent prestack, if  $\forall n \geq 0 \exists Z_n \in \text{Sch}^{\text{aff}}$ ,  
then  $Z \in \text{Sch}$ .

We omit the discussion of finiteness conditions for the moment and just define:

a morphism  $f: X \rightarrow Y$  of prestacks is schematic if  $\forall S \in \text{Sch}^{\text{aff}}$   
&  $S \rightarrow Y$  the pullback  $X \times_Y S \in \text{Sch}$ .

Clearly, one can prove:

Lemma: For  $X \rightarrow Y$  a schematic map, if  $Z \in \text{Sch}$  then  $X \times_Y Z \in \text{Sch}$ .